Concurrent Models of Computation

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EECS 219D: Concurrent Models of Computation
Fall 2011
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Week 1: Process Networks

Logistics
Class web page:
http://embedded.eecs.berkeley.edu/concurrency

Project
Paper (and paper review)
Homework
Reading
Technology:
  ● Ptolemy II
  ● Java
  ● Eclipse
  ● LaTeX
Homework

- Issued roughly every two weeks
- Will leverage a common technology base:
  - Java
  - Eclipse
  - Ptolemy II
- First assignment is on the web

Project

- Conference (workshop) paper quality expected
- Papers will be “submitted” and “reviewed” by you
- Presentations will be workshop like
- Teams up to two are encouraged
- Leveraging the technology base is encouraged
- Many project suggestions are on the web
  - In almost all cases, I have a fairly clear idea of how to start. Come talk to me if one of these looks interesting
Introduction to Edward A. Lee

- Working in embedded software since 1978, when I was writing assembly code for 8-bit microcontrollers to control biomedical robotic machines. From 1980-82, I was writing assembly code for the AT&T DSP-1 to implement modems at Bell Labs.
- BS ’79 (Yale, Double major: CS and EE)
  SM ’81 (MIT, EECS)
  PhD ’86 (Berkeley, EECS)
- Berkeley EECS faculty since 1986
- One of four directors of Chess, the Berkeley Center for Hybrid and Embedded Software Systems
- Director of the Berkeley Ptolemy project
- Co-author of nine books (on embedded systems, digital communications, signals and systems, dataflow)
- Chair of EE, then EECS, from Jan. 2005- June 2008.
- Co-founder of BDTI, Inc., a 19 year old technology company

Who are you?
Model of Computation

NIST:
“A formal, abstract definition of a computer.”
Examples: Turing machine, random access machine, primitive recursive, cellular automaton, finite state machine, ...

Wikipedia (on 1/18/09):
“the definition of the set of allowable operations used in computation and their respective costs.”

“In model-driven engineering, the model of computation explains how the behaviour of the whole system is the result of the behaviour of each of its components.”
Concurency

From the Latin, concurrere, “run together”

Discussion:
Is concurrency hard?
Potential Confusion

- Concurrent vs. parallel
- Concurrent vs. nondeterminate

Kahn Process Networks (PN)

A Concurrent Model of Computation (MoC)

- A set of components called actors.
- Each representing a sequential procedure.
- Where steps in these procedures receive or send messages to other actors (or perform local operations).
- Messages are communicated asynchronously with unbounded buffers.
- A procedure can always send a message. It does not need to wait for the recipient to be ready to receive.
- Messages are delivered reliably and in order.
- When a procedure attempts to receive a message, that attempt blocks the procedure until a message is available.
Coarse History

- Semantics given by Gilles Kahn in 1974.
  - Fixed points of continuous and monotonic functions
- More limited form given by Kahn and MacQueen in 1977.
  - Blocking reads and nonblocking writes.
- Generalizations to nondeterministic systems
  - Kosinski [1978], Stark [1980s], …
- Bounded memory execution given by Parks in 1995.
  - Solves an undecidable problem.
- Debate over validity of this policy, Geilen and Basten 2003.
  - Relationship between denotational and operational semantics.
- Many related models intertwined.
  - Actors (Hewitt, Agha), CSP (Hoare), CCS (Milner), Interaction (Wegner), Streams (Broy, …), Dataflow (Dennis, Arvind, …)...
Question 1:
Is “Fair” Thread Scheduling a Good Idea?

In the following model, what happens if every thread is given an equal opportunity to run?
Rendezvous: An Alternative Communication Mechanism with Bounded Buffers

Rendezvous underlies CSP (Hoare), CCS (Milner), and Statecharts (Harel)

```java
while(true) {
    data1 = in1.get();
    data2 = in2.get();
    ... do something with it ...
}
```

Discussion: What should a fork do?

Processes must both reach this point before either can continue.

Discussion

How does this program compare under rendezvous communication vs. process networks?
Question 2: Should we use Rendezvous Here?

The control signal now depends on the source data.

A Practical Application with this Structure

Consider collecting time-stamped trades from commodities markets around the world and merging them into a single time-stamped stream. The CONTROL actors could compare time stamps, with logic like this:

```java
data1 = topPort.get();
data2 = bottomPort.get();
while (true) {
    if (data1.time < data2.time) {
        output.send(true);
        data1 = topPort.get();
    } else {
        output.send(false);
        data2 = bottomPort.get();
    }
}
```
Question 3: How about Demand-Driven (Lazy) Execution?

In demand-driven execution, a process is stalled unless its outputs are required by a downstream process.

Will Demand-Driven Execution Work Here?
Question 4: Will Data-Driven Execution Work?

In data-driven execution, a process is stalled unless it has input data. What about the processes with no inputs?

Things are not looking good…

We have ruled out:
- Fair execution.
- Rendezvous communication.
- Demand-driven execution.
- Data-driven execution.

For all the examples given so far, there is an obvious execution policy that does what we want. Is there a general policy that will always deliver that obvious policy?

Are there models for which the policy is not so obvious?
Question 5:
What is the “Correct” Execution of This Model?

```
while(true) {
  data1 = in1.get();
  data2 = in2.get();
  ... do something with it ...
}
```

Question 6:
What is the Correct Behavior of this Model?
Question 7: How to support nondeterminism?

Merging of streams is needed for some applications. Does this require fairness? What does fairness mean?

Properties of PN (Two Big Topics)

- Assuming “well-behaved” actors, a PN network is determinate in that the sequence of tokens on each arc is independent of the thread scheduling strategy.
  - Making this statement precise, however, is nontrivial.

- PN is Turing complete.
  - Given only boolean tokens, memoryless functional actors, Switch, Select, and initial tokens, one can implement a universal Turing machine.
  - Whether a PN network deadlocks is undecidable.
  - Whether buffers grow without bound is undecidable.
Where This is Going

A signal is a sequence of values

Define a prefix order:

\( a \sqsubseteq a' \)

means that \( a \) is a prefix of \( a' \).

Actors are **monotonic** functions:

\[
\begin{align*}
    a \sqsubseteq a' & \implies f(a) \sqsubseteq f(a')
\end{align*}
\]

Stronger condition: Actors are **continuous** functions (intuitively: they don’t wait forever to produce outputs).

PN Semantics of Composition (Kahn, ’74)

This Approach to Semantics is “Tarskian”

If the components are deterministic, the composition is deterministic.

\[
\begin{align*}
x = y & \implies f(x) = x
\end{align*}
\]

Fixed point theorem:

- Continuous function has a unique least fixed point
- Execution procedure for finding that fixed point
- Successive approximations to the fixed point
What is Order?

Intuition:
1. $0 < 1$
2. $1 < \infty$
3. child < parent
4. child > parent
5. $11,000/3,501$ is a better approximation to $\pi$ than $22/7$
6. integer $n$ is a divisor of integer $m$.
7. Set $A$ is a subset of set $B$.

Which of these are partial orders?

Relations
- A relation $R$ from $A$ to $B$ is a subset of $A \times B$
- A function $F$ from $A$ to $B$ is a relation where $(a, b) \in R$ and $(a, b') \in R \Rightarrow b = b'$
- A binary relation $R$ on $A$ is a subset of $A \times A$
- A binary relation $R$ on $A$ is reflexive if $\forall a \in A$, $(a, a) \in R$
- A binary relation $R$ on $A$ is symmetric if $(a, b) \in R \Rightarrow (b, a) \in R$
- A binary relation $R$ on $A$ is antisymmetric if $(a, b) \in R$ and $(b, a) \in R \Rightarrow a = b$
- A binary relation $R$ on $A$ is transitive if $(a, b) \in R$ and $(b, c) \in R \Rightarrow (a, c) \in R$
Infix Notation for Binary Relations

- \((a, b) \in R\) can be written \(a R b\)
- A symbol can be used instead of \(R\). For examples:
  - \(\leq \subseteq \mathbb{N} \times \mathbb{N}\) is a relation.
  - \((a, b) \in \leq\) is written \(a \leq b\)
- A function \(f \in (A, B)\) can be written \(f : A \rightarrow B\)

Partial Orders

A **partial order** on the set \(A\) is a binary relation \(\leq\) that is:

For all \(a, b, c \in A\),
- reflexive: \(a \leq a\)
- antisymmetric: \(a \leq b\) and \(b \leq a \Rightarrow a = b\)
- transitive: \(a \leq b\) and \(b \leq c \Rightarrow a \leq c\)

A **partially ordered set** (poset) is a set \(A\) and a binary relation \(\leq\), written \((A, \leq)\).
Strict Partial Order

For every partial order \( \leq \) there is a strict partial order \( < \) where \( a < b \) if and only if \( a \leq b \) and \( a \neq b \).

A strict poset is a set and a strict partial order.

Total Orders

Elements \( a \) and \( b \) of a poset \((A, \leq)\) are comparable if either \( a \leq b \) or \( b \leq a \). Otherwise they are incomparable.

A poset \((A, \leq)\) is totally ordered if every pair of elements is comparable.

Totally ordered sets are also called linearly ordered sets or chains.
Quiz

1. Is the set of integers with the usual numerical ordering a well-ordered set? (A well-ordered set is a set where every non-empty subset has a least element.)

2. Given a set $A$ and its powerset (set of all subsets) $P(A)$, is $(P(A), \subseteq)$ a poset? A chain?

3. For $A = \{a, b, c\}$ (a set of three letters), find a well-ordered subset of $(P(A), \subseteq)$.

Answers

1. Is the set of integers with the usual numerical ordering a well-ordered set?
   No. The set itself is a chain with no least element.

2. Given a set $A$ and its powerset (set of all subsets) $P(A)$, is $(P(A), \subseteq)$ a poset? A chain?
   It is a poset, but not a chain.

3. For $A = \{a, b, c\}$ (a set of three letters), find a well-ordered subset of $(P(A), \subseteq)$.
   One possibility: \{∅, \{a\}, \{a, b\}, \{a, b, c\}\}
Pertinent Example: Prefix Orders

Let $T$ be a type (a set of values).
Let $T^\ast$ be the set of all finite and infinite sequences of elements of $T$, including the empty sequence $\bot$ (bottom).

Let $\sqsubseteq$ be a binary relation on $T^\ast$ such that $a \sqsubseteq b$ if $a$ is a prefix of $b$. That is, for all $n$ in $N$ such that $a(n)$ is defined, then $b(n)$ is defined and $a(n) = b(n)$.

This is called a prefix order.

During execution, the outputs of a PN actor form a well-ordered subset of $(T^\ast, \sqsubseteq)$.

Summary

- Concurrent models of computation
- Process networks as an example
- Intuitive model, but many subtle corner cases
- Need a solid theory underlying it
- Posets

- Next time:
  - give meaning to all programs
  - develop an execution policy